CSCI 204: Data Structures \& Algorithms

## Recursion 2

## Recursive binary search

- Idea:
- Compare the target with the list item in the middle
- If found, stop;
- If target is greater than the middle, search the second half, otherwise, search the first half
- Base case(s):
- Found or the list is exhausted
- Recursive case:
- Search the first half
- Search the second half


## Recursive binary search

if left > right: \# not found
if left > right: \# not found
id = (Left + right) // 4 N
id = (Left + right) // 4 N
Mif rens |mid] < target: \# search for upper half
Mif rens |mid] < target: \# search for upper half
Leturn bin_search(rumus, target, left, right)
Leturn bin_search(rumus, target, left, right)


manmen
manmen
Ma,m,
Ma,m,

## List all permutations

- Make every element in the list as a prefix, one at a time, do it recursively
- E.g., 'abcd'
- 'a' + recursively('bcd')
- 'b' + recursively('acd')
$-{ }^{\prime} c^{\prime}+$ recursively('abd')
- 'd' + recursively('abc')

Now let's do the workshop.

## Check if a number is a prime

- Ideas: to determine if $b$ is a prime, we check if $b \% x$ == 0 (divisible) consecutively ...
- E.g., 5 : we check $5 \% 4,5 \% 3,5 \% 2,5 \% 1$, when x reaches 1 , we know 5 is a prime
- E.g., 6 : we check $6 \% 5,6 \% 4,6 \% 3$ which is 0 , stop, 6 is not a prime
- The task is to place 8 queens onto a chessboard such that no queen can attack another queen.
- Uses a standard $8 \times 8$ chess board.
- There are 92 solutions to this problem.


## Queen's Moves

- The queen can move and attack any piece of the opponent by moving in any direction along a straight line.



## Sample Solutions



## 4-Queens Problem

- To develop an algorithm, we consider the smaller 4-queens problem.
- Since no two queens can occupy the same column, we can proceed one column at a time.
- Place a queen in position ( 0,0 ).



## 4-Queens Problem

- We move to the second column and
place a queen at position $(2,1)$

| 畨 | $x$ | $x$ | $x$ |
| :---: | :---: | :---: | :---: |
| $x$ | $x$ | $x$ |  |
| $x$ | $\omega$ | $x$ |  |
| $x$ | $x$ | $x$ | $x$ |

## 4-Queens Problem

- This move eliminates a number of squares for the placement of additional queens.

| 畨 | x | x | x |
| :---: | :---: | :---: | :---: |
| x | x |  |  |
| x |  | x |  |
| x |  |  | x |

## 4-Queens Problem

- The $3^{\text {rd }}$ queen should be placed in the $3^{\text {rd }}$ column.
- But there are no open cells in the third column.
- So there is no solution based on the placement of the first 2 queens.

| W. | $x$ | $x$ | $x$ |
| :---: | :---: | :---: | :---: |
| $x$ | $x$ | $x$ |  |
| $x$ | $\boldsymbol{w}$ | $x$ |  |
| $x$ | $x$ | $x$ | $x$ |

## 4－Queens Problem

－We have to backtrack：
－go back to the previous column
－pickup the last queen placed
－try to find another valid cell in that column．

| 畨 | $x$ | $x$ | $x$ |
| :---: | :---: | :---: | :---: |
| $x$ | $x$ |  |  |
| $x$ |  | $x$ |  |
| $x$ |  |  | $x$ |

## 4－Queens Problem

－In the $3^{\text {rd }}$ column，we can now place a queen at position（1，2）．
－But now we have no open slots in the $4^{\text {th }}$ column．

| 类 | x | x | x |
| :---: | :---: | :---: | :---: |
| x | x | $\boldsymbol{\omega}$ | x |
| x | x | x | x |
| x | U | x | x |

## 4－Queens Problem

－Place a queen at position $(3,1)$ and move forward．

| $\boldsymbol{\omega}$ | x | x | x |
| :---: | :---: | :---: | :---: |
| x | x |  |  |
| x | x | x |  |
| x | $\boldsymbol{\omega}$ | x | x |

## 4－Queens Problem

－We again must backtrack and pick up the queen from the $3^{\text {rd }}$ column．
－But there are no other empty cells in the $3^{\text {rd }}$ column．

| 业 | $x$ | $x$ | $x$ |
| :---: | :---: | :---: | :---: |
| $x$ | $x$ | $x$ |  |
| $x$ | $x$ | $x$ |  |
| $x$ | 震 | $x$ | $x$ |

## 4－Queens Problem

－We must backtrack yet again and pick up the queen from the $2^{\text {rd }}$ column．
－But there are no other empty cells in the $2^{\text {nd }}$ column either．

| wiw | $x$ | $x$ | $x$ |
| :---: | :---: | :---: | :---: |
| $x$ | $x$ |  |  |
| $x$ | $x$ | $x$ |  |
| $x$ | $x$ |  | $x$ |

## 4－Queens Problem

－So we backtrack one more time and pick up the queen from the $1^{\text {st }}$ column．
－We then try again to place a queen in the $1^{\text {st }}$ column．


## 4-Queens Problem

- In the $1^{\text {st }}$ column, we can place a queen at position ( 1,0 ).

| $x$ | $x$ |  |  |
| :---: | :---: | :---: | :---: |
| 䒼 | $x$ | $x$ | $x$ |
| $x$ | $x$ |  |  |
| $x$ |  | $x$ |  |

## N-Queens Board ADT

- The $n$-queens board is used for positioning queens on a square board for use in solving the $n$-queens problem.
- consists of $n \times n$ squares.
- each square is identified by index [0...n)

| - NQueensBoard( n ) | - placeQueen( row, col ) |
| :---: | :---: |
| - size() | - removeQueen(row, |
| - numQueens() | - reset() |
| - unguarded( row, col ) | - draw() |

## 4-Queens Problem

- We again continue with the process and attempt to find open positions in each of the remaining columns.
- We can use a similar approach to solve the original 8-queens problem.



## 8-Queens Solution

def solveNQueens ( board, col ):
if board.numQueens() == board.size() :
return True
else :
for row in range( board.size() ):
if board. unguarded ( row, col ):
board.placeQueen ( row, col )
if board. solveNQueens (board, col+1 )
return True
else
board.removeQueen ( row, col )
return False

